



STIC Search Report

EIC 2100

STIC Database Tracking Number: 143598

TO: Thu Ha Nguyen
Location: 4A71
Art Unit : 2155
Friday, January 28, 2005

Case Serial Number: 09905080

From: David Holloway
Location: EIC 2100
RND 4B19
Phone: 2-3528

david.holloway@uspto.gov

Search Notes

Dear Examiner Nguyen,

Attached please find your search results for above-referenced case.
Please contact me if you have any questions or would like a re-focused search.

David

Set	Items	Description
S1	6396310	BANDWIDTH? OR FLOW? OR RESOURCE? OR LOAD? OR CAPACITY OR D-ATA() RATE?
S2	374665	S1(3N) (CONTROL? OR BALANC? OR MANAGE? OR ADMINIST?)
S3	1629690	SWITCH? OR NODE? OR SERVER? OR ROUTER? OR BRIDGE? OR GATEWAY?
S4	8004	(VIRTUAL? OR LOGICAL?) (2N) (LANE? OR CHANNEL? OR ROUTE OR ROUTES OR PATH OR PATHS OR THREAD?)
S5	440245	QUEU? OR CACH? OR BUFFER? OR TEMPORAR?() (STORAGE? OR MEMORY?)
S6	4663	INFINIBAND? OR HCA
S7	691	S2 AND S3 AND S4
S8	194	S7 AND S5
S9	0	S8 AND S6
S10	1	S7 AND S6
S11	38	S2 AND S6
S12	44	S2(5N)S3 AND S4 AND S5
S13	62	S8 AND (NOTIF? OR PING OR ACK OR ACKNOWLEDG? OR MESSAG?)
S14	125	S11 OR S12 OR S13
S15	84	RD (unique items)
S16	59	S15 NOT PY>2000
S17	33	S16 AND S13
S18	34	S10 OR S17
File	8:Ei Compendex(R) 1970-2005/Jan W3	
	(c) 2005 Elsevier Eng. Info. Inc.	
File	35:Dissertation Abs Online 1861-2004/Dec	
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File	65:Inside Conferences 1993-2005/Jan W4	
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File	2:INSPEC 1969-2005/Jan W3	
	(c) 2005 Institution of Electrical Engineers	
File	94:JICST-EPlus 1985-2005/Dec W3	
	(c)2005 Japan Science and Tech Corp(JST)	
File	111:TGG Natl.Newspaper Index(SM) 1979-2005/Jan 25	
	(c) 2005 The Gale Group	
File	6:NTIS 1964-2005/Jan W3	
	(c) 2005 NTIS, Intl Copyrgh All Rights Res	
File	144:Pascal 1973-2005/Jan W3	
	(c) 2005 INIST/CNRS	
File	434:SciSearch(R) Cited Ref Sci 1974-1989/Dec	
	(c) 1998 Inst for Sci Info	
File	34:SciSearch(R) Cited Ref Sci 1990-2005/Jan W4	
	(c) 2005 Inst for Sci Info	
File	62:SPIN(R) 1975-2005/Nov W1	
	(c) 2005 American Institute of Physics	
File	99:Wilson Appl. Sci & Tech Abs 1983-2004/Nov	
	(c) 2004 The HW Wilson Co.	
File	95:TEME-Technology & Management 1989-2004/Jun W1	
	(c) 2004 FIZ TECHNIK	

18/5/4 (Item 4 from file: 8)
DIALOG(R)File 8:Ei Compendex(R)
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04643528 E.I. No: EIP97033564128

Title: Analysis and optimization of pacing window flow control with admission delay

Author: Suk, Jung-Bong; Cassandras, Christos G.

Corporate Source: Yonsei Univ

Source: IEICE Transactions on Information and Systems v E79-D n 12 Dec 1996. p 1663-1675

Publication Year: 1996

CODEN: ITISEF ISSN: 0916-8532

Language: English

Document Type: JA; (Journal Article) Treatment: T; (Theoretical)

Journal Announcement: 9704W5

Abstract: This paper provides a queueing model analysis of virtual route networks for which a pacing window flow control mechanism is employed with an input queue included. The input queue is introduced into the model to describe the waiting system where messages prevented from entering the network are stored in first-come first-serve manner. Both cases of finite and infinite capacity are considered. The model leads to a Markovian queueing system, which is fully solved through appropriate use of matrix-geometric methods. The empirical rule is that the optimum window size which maximizes the power criterion including the admission delay is nearly twice the number of hops (nodes of the network). Simulations are presented to verify the analytical results. Finally, performance comparisons with the sliding window protocol are made. Our results show that although the average number of messages in the network is higher for the pacing window case, when the input queue delay is taken into consideration the overall performance of the pacing window protocol is better than that of the sliding window. (Author abstract) 11 Refs.

Descriptors: *Congestion control (communication); Computer networks; Data communication systems; Optimization; Mathematical models; Data storage equipment; Markov processes; Matrix algebra; Computer simulation; Network protocols

Identifiers: Pacing window flow control ; Markovian queueing model; Virtual route networks; Admission delay; Matrix geometric methods

Classification Codes:

716.1 (Information & Communication Theory); 921.5 (Optimization Techniques); 921.6 (Numerical Methods); 722.1 (Data Storage, Equipment & Techniques); 922.1 (Probability Theory); 921.1 (Algebra)
716 (Radar, Radio & TV Electronic Equipment); 921 (Applied Mathematics); 722 (Computer Hardware); 922 (Statistical Methods)
71 (ELECTRONICS & COMMUNICATIONS); 92 (ENGINEERING MATHEMATICS); 72 (COMPUTERS & DATA PROCESSING)

18/5/6 (Item 6 from file: 8)
DIALOG(R)File 8:EI Compendex(R)
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04114530 E.I. No: EIP95032625179

Title: Modeling virtual channel flow control in hypercubes

Author: Boura, Younes M.; Das, Chita R.

Corporate Source: Pennsylvania State Univ, University Park, PA, USA

Conference Title: Proceedings of the 1st IEEE Symposium on
High-Performance Computer Architecture

Conference Location: Raleigh, NC, USA Conference Date:
19950122-19950125

Sponsor: IEEE

E.I. Conference No.: 42661

Source: IEEE High-Performance Computer Architecture Symposium Proceedings
1995. IEEE, Los Alamitos, CA, USA, 95TH8026. p 166-175

Publication Year: 1995

CODEN: 001971

Language: English

Document Type: CA; (Conference Article) Treatment: G; (General Review);
T; (Theoretical)

Journal Announcement: 9505W3

Abstract: An analytical model for virtual channel flow control in n-dimensional hypercubes using the e-cube routing algorithm is developed. The model is based on determining the values of the different components that make up the average message latency. These components include the message transfer time, the blocking delay at each dimension, the multiplexing delay at each dimension, and the waiting delay at the source node. The first two components are determined using a probabilistic analysis. The average degree of multiplexing is determined using a Markov model, and the waiting delay at the source node is determined using an M/M/m queueing system. The model is fairly accurate in predicting the average message latency for different message sizes and a varying number of virtual channels per physical channel. 13 Refs.

Descriptors: *Interconnection networks; Mathematical models; Algorithms; Markov processes; Process control; Queueing theory; Probability; Switching ; Multiplexing; Data transfer

Identifiers: Virtual channel flow control ; Analytical model; E cube routing algorithm; N dimensional hypercubes; Message latency; Markov model; Wormhole switching ; M m m queueing system; Probabilistic analysis; Multiplexing delay

Classification Codes:

721.3 (Computer Circuits); 921.6 (Numerical Methods); 731.3 (Specific Variables Control); 922.1 (Probability Theory); 723.2 (Data Processing); 722.3 (Data Communication, Equipment & Techniques)

721 (Computer Circuits & Logic Elements); 921 (Applied Mathematics); 731 (Automatic Control Principles); 922 (Statistical Methods); 723 (Computer Software); 722 (Computer Hardware)

72 (COMPUTERS & DATA PROCESSING); 92 (ENGINEERING MATHEMATICS); 73 (CONTROL ENGINEERING)

18/5/11 (Item 1 from file: 35)
DIALOG(R)File 35:Dissertation Abs Online
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01687281 ORDER NO: NOT AVAILABLE FROM UNIVERSITY MICROFILMS INT'L.
RESOURCE MANAGEMENT FOR CLASSES OF SERVICE IN WORMHOLE NETWORKS
(QUALITY OF SERVICE, ROUTING)
Author: SMAI, ABDEL-HALIM
Degree: TEKN.DR
Year: 1997
Corporate Source/Institution: KUNGLIGA TEKNISKA HOGSKOLAN (SWEDEN) (1022)
Source: VOLUME 60/02-C OF DISSERTATION ABSTRACTS INTERNATIONAL.
PAGE 448. 142 PAGES
Descriptors: COMPUTER SCIENCE ; ENGINEERING, SYSTEM SCIENCE
Descriptor Codes: 0984; 0790
Publisher: ROYAL INSTITUTE OF TECHNOLOGY, S-100 44 STOCKHOLM, SWEDEN

Parallel computers can provide high compute, storage and communication rates for scientific, engineering and multimedia applications. Nevertheless, designing an efficient communication subsystem for a parallel computer continues to be a challenging problem. For instance, existing multicomputer networks are often not adapted to support classes of service. Network routers use arbitration schemes such as first-come-first-served and round-robin which are not suitable to deliver such a service diversity.

Due to its simplicity, low cost and high performance, the wormhole routing switching technique has become the method of choice in building high-speed networks today. However, the interaction between multiple messages competing for shared resources and the blocking property of wormhole routing increase the risk for network congestion and make it difficult to support good quality of service.

In this thesis we propose and evaluate resource management techniques to support special services for particular classes of traffic in wormhole networks. We address issues related to management of network resources such as point-to-point link bandwidth and buffer storage, and the problem of congestion. Evaluation is done by simulation.

First, we propose a low-latency virtual channel flow control for prioritized traffic, called Prioritized Demand Multiplexing (PDM). Two schemes are derived, strict and relaxed PDM, which are based both on round-robin and priority scheduling. Relaxed PDM is particularly appropriate to support multiple number of classes of traffic, based on the concept of extra chances. Based on PDM, we propose fast absorb flow control and distance-based flow control. The distance-based approach relies on a distance-based priority mapping scheme, and the fast absorb flow control on priority setting at the destination instead of the source.

We introduce grouping of virtual channels to provide support for classes of service. Two approaches are proposed: static and semi-static. In the semi-static approach, messages are allowed to borrow virtual channels reserved for messages with lower priority. Further, we propose dynamic configuration of virtual channels.

Finally, a general approach to global reactive congestion control is developed. The approach uses timeout mechanism to detect congestion, and exploits control lines such as those used for handshaking in the flit-level flow control of wormhole routers to distribute information about congestion. It is also based on a throttling mechanism that limits the demands placed by the network interface and the processing element.

18/5/19 (Item 2 from file: 2)

DIALOG(R)File 2:INSPEC

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5836360 INSPEC Abstract Number: C9803-5220P-066

Title: **Ghost packets: a deadlock-free solution for k-ary n-cube networks**

Author(s): Carrion, C.; Izu, C.; Gregorio, J.A.; Vallejo, F.; Beivide, R.

Author Affiliation: Dept. de Electron., Cantabria Univ., Santander, Spain

Conference Title: Proceedings of the Sixth Euromicro Workshop on Parallel and Distributed Processing - PDP'98 - (Cat. No. 98EX134) p.133-9

Publisher: IEEE Comput. Soc, Los Alamitos, CA, USA

Publication Date: 1998 Country of Publication: USA xiii+520 pp.

ISBN: 0 8186 8332 5 Material Identity Number: XX98-00234

U.S. Copyright Clearance Center Code: 0 8186 8332 5/98/\$10.00

Conference Title: Proceedings of the Sixth Euromicro Workshop on Parallel and Distributed Processing - PDP '98 -

Conference Sponsor: Dept. Electron. Univ. York; Univ. Complutense Madrid

Conference Date: 21-23 Jan. 1998 Conference Location: Madrid, Spain

Language: English Document Type: Conference Paper (PA)

Treatment: Practical (P)

Abstract: Improving interconnection subsystems is crucial for the overall performance of a multicomputer system. Hence, a theoretical presentation of a new deadlockfree **message** flow model for k-ary n-cube networks is developed in this paper. The key idea of this **flow control** mechanism is to preserve enough free resources for each possible routing dependency cycle, so that packet progress will be guaranteed. Based on this algorithm, we have proposed a simple **router** structure for a 2-ary n-cube topology with dimensional order routing. Edge or shared **buffering** can be used, requiring a minimum capacity of one packet per **channel**. **Virtual channels** are eliminated, reducing **router** complexity and, consequently, decreasing network latency at lour loads. In fact, the performance evaluation for the 2-ary n-cube with different loads shows an improvement in the latency parameter of about 20% with respect to a deterministic routing with two **virtual channels**. (14 Refs)

Subfile: C

Descriptors: computational complexity; concurrency control; multiprocessor interconnection networks; performance evaluation; system recovery

Identifiers: ghost packets; deadlock-free solution; k-ary n-cube networks ; interconnection subsystems; performance; **message** flow model; **virtual channels** ; complexity; deterministic routing

Class Codes: C5220P (Parallel architecture); C5470 (Performance evaluation and testing); C4230M (Multiprocessor interconnection); C4240C (Computational complexity); C5440 (Multiprocessing systems); C6150G (Diagnostic, testing, debugging and evaluating systems)

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Set	Items	Description
S1	2742141	BANDWIDTH? OR FLOW? OR RESOURCE? OR LOAD? OR CAPACITY OR D- ATA() RATE?
S2	214150	S1(3N) (CONTROL? OR BALANC? OR MANAGE? OR ADMINIST?)
S3	1625485	SWITCH? OR NODE? OR SERVER? OR ROUTER? OR BRIDGE? OR GATEW- AY?
S4	4610	(VIRTUAL? OR LOGICAL?) (2N) (LANE? OR CHANNEL? OR ROUTE OR R- OUTES OR PATH OR PATHS OR THREAD?)
S5	294581	QUEU? OR CACH? OR BUFFER? OR TEMPORAR? () (STORAGE? OR MEMOR- ?)
S6	347	INFINIBAND? OR HCA
S7	140	S2 AND S3 AND S4
S8	32	S7 AND S5
S9	0	S8 AND S6
S10	0	S7 AND S6
S11	13	S2 AND S6
S12	45	S8 OR S11
S13	29	S12 NOT AD>20010718

File 347:JAPIO Nov 1976-2004/Aug(Updated 041203)
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File 350:Derwent WPIX 1963-2005/UD,UM &UP=200504
(c) 2005 Thomson Derwent

13/5/2 (Item 2 from file: 347)
DIALOG(R) File 347:JAPIO
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06552957 **Image available**
CONTROL PATH BAND WARRANT SYSTEM IN HIGH SPEED ROUTER

PUB. NO.: 2000-138686 [JP 2000138686 A]
PUBLISHED: May 16, 2000 (20000516)
INVENTOR(s): OCHIAI TAMIYA
YASUDA HIROKAZU
APPLICANT(s): TOSHIBA CORP
TOSHIBA TELECOMMUNICATION SYSTEM ENGINEERING CORP
APPL. NO.: 10-311319 [JP 98311319]
FILED: October 30, 1998 (19981030)
INTL CLASS: H04L-012/28; H04Q-003/00

ABSTRACT

PROBLEM TO BE SOLVED: To attain high speed transfer by reserving a path for a control flow so as to maintain a cut-through path even in the case of excess traffic.

SOLUTION: An output side of an asynchronous transfer mode ATM switch section 20 of a high speed router discriminates packets being output objects in the unit of virtual channel VC, a control packet is stored in a priority buffer 21a so as to give priority to a control path through which the control packet passes according to the discrimination result and outputted with priority thereby reserving a communication band of the control path. Thus, even on the occurrence of excess traffic, lack of communication of the control packet is avoided so as to maintain a cut-through path and to attain high speed transfer.

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13/5/5 (Item 5 from file: 347)
DIALOG(R) File 347:JAPIO
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05646913 **Image available**

FLOW CONTROL METHOD AND MOBILE COMMUNICATION NETWORK HAVING FLOW
CONTROL FUNCTION

PUB. NO.: 09-261713 [JP 9261713 A]
PUBLISHED: October 03, 1997 (19971003)

INVENTOR(s): NISHIO MASAYA
SHINAGAWA NORITERU
WATANABE YONEO
TANAKA MOTOHARU

APPLICANT(s): Y R P IDO TSUSHIN KIBAN GIJUTSU KENKYUSHO KK [000000] (A
Japanese Company or Corporation), JP (Japan)

APPL. NO.: 08-093018 [JP 9693018]

FILED: March 25, 1996 (19960325)

INTL CLASS: [6] H04Q-007/22; H04L-012/28; H04Q-003/00; H04Q-007/28

JAPIO CLASS: 44.2 (COMMUNICATION -- Transmission Systems); 44.3
(COMMUNICATION -- Telegraphy); 44.4 (COMMUNICATION --
Telephone)

ABSTRACT

PROBLEM TO BE SOLVED: To provide a flow control method and flow
controller between ATM nodes for hand-over in an ATM mobile
communication network in which a data loss is eliminated with a small
memory capacity.

SOLUTION: A mobile station 11 makes hand-over from a radio base station 21
to a radio base station 22 attended with its movement, and the radio base
station 21 sends a hand-over start notice cell to an in-zone control
station 31. A cell switch controller 81 of the in-zone control station 31
selects a corresponding ATM virtual channel and 'holds' output port
information relating to the channel and a flow control module 103
stores a corresponding cell to an output hold queue 113. At the end of
hand-over, the radio base station 22 sends a hand-over end notice cell, the
cell switch controller 91 releases holding of the output port information
and a cell switch 81 outputs a cell from the output hold queue .

13/5/6 (Item 6 from file: 347)
DIALOG(R) File 347:JAPIO
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02425643 **Image available**
PACKET FLOW CONTROLLING SYSTEM

PUB. NO.: 63-042543 [JP 63042543 A]
PUBLISHED: February 23, 1988 (19880223)
INVENTOR(s): ONISHI KOICHI
NOGUCHI KIYOHIRO
APPLICANT(s): NIPPON TELEGR & TELEPH CORP <NTT> [000422] (A Japanese
Company or Corporation), JP (Japan)
APPL. NO.: 61-186297 [JP 86186297]
FILED: August 08, 1986 (19860808)
INTL CLASS: [4] H04L-011/20
JAPIO CLASS: 44.3 (COMMUNICATION -- Telegraphy)
JOURNAL: Section: E, Section No. 635, Vol. 12, No. 256, Pg. 33, July
19, 1988 (19880719)

ABSTRACT

PURPOSE: To realize the **flow control** which can cope with the overload condition of traffic, by controlling the transmission throughput by a packet terminal so that its own transmission throughput does not exceed a reported maximum throughput and abandoning excess packets or disconnecting a logical channel by an exchange if the packet terminal transmits packets with a throughput exceeding the reported value.

CONSTITUTION: The number of packets is counted by a transmission packet (throughput) counter 14 in every throughput decision cycle which is a certain period reported by a throughput cycle switching report line 51, and the counted value is defined as the transmission throughput in the current cycle and is arranged with a network at the time of originating a call or the like, and it is checked by a throughput value comparator 16 whether this throughput value exceeds a maximum transmission throughput value stored in a throughput value holding device 15 or not. If it does not exceed, a transmission permission report line 54 is set to the transmittable state and a packet transmission (throughput) controller 12 transmits packets from a transmission packet buffer 13. If it exceeds, the transmission permission report line 54 is set to the untransmittable state and the packet transmission (throughput) controller 12 stops the transmission.

13/5/8 (Item 2 from file: 350)

DIALOG(R) File 350:Derwent WPIX

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015050340 **Image available**

WPI Acc No: 2003-110856/200310

XRPX Acc No: N03-088171

Link packet scheduler for use with LAN, WAN, compares accumulated free credits and current buffer receive utilization with programmable credit and utilization threshold respectively.

Patent Assignee: REOHR R D (REOH-I); SUSNOW D S (SUSN-I)

Inventor: REOHR R D; SUSNOW D S

Number of Countries: 001 Number of Patents: 001

Patent Family:

Patent No	Kind	Date	Applicat No	Kind	Date	Week
US 20020159385	A1	20021031	US 2001842019	A	20010426	200310 B

Priority Applications (No Type Date): US 2001842019 A 20010426

Patent Details:

Patent No	Kind	Lan Pg	Main IPC	Filing Notes
US 20020159385	A1	21	H04J-003/14	

Abstract (Basic): US 20020159385 A1

NOVELTY - An N-bit counter accumulates free credits relinquished when a data packet is removed or when a link packet is received. The comparators compare the accumulated credits and a current buffer receive utilization with a programmable credit and utilization threshold respectively. A logic device monitors the comparators and schedule link packets accordingly.

DETAILED DESCRIPTION - INDEPENDENT CLAIMS are included for the following:

- (1) Data network; and
- (2) Method of flow control of a link packet in a host-fabric adapter.

USE - Link packet scheduler for use with data network (claimed) such as LAN, WAN, CAN, MAN, global area network, wireless personal network, system area network including network using next generation input/output, future input/output, Infiniband, server net, etc., input/output hardware adapters and chip sets.

ADVANTAGE - The utilization of credit based flow control prevents the transmitter from sending data packets unless the receiver has room for accepting the packets. Thus loss of data packets due to receive buffer overflow is prevented. Programmable credit acquisition threshold enables credits to be returned as they become available not every 65,536 symbol times which results in a significant network performance increase.

DESCRIPTION OF DRAWING(S) - The figure shows the explanatory view of the link packet flow control mechanism.

pp; 21 DwgNo 7/9

Title Terms: LINK; PACKET; LAN; WAN; COMPARE; ACCUMULATE; FREE; CREDIT; CURRENT; BUFFER; RECEIVE; UTILISE; PROGRAM; CREDIT; UTILISE; THRESHOLD; RESPECTIVE

Derwent Class: T01; W01

International Patent Class (Main): H04J-003/14

International Patent Class (Additional): H04J-001/16

File Segment: EPI

13/5/10 (Item 4 from file: 350)
DIALOG(R) File 350:Derwent WPIX
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014834841 **Image available**
WPI Acc No: 2002-655547/200270
XRPX Acc No: N02-518032

Buffering method of packets transmitted to infiniband port, involves transmitting flow control credits to resource buffer device to stop transmission of packets when memory is filled with specific amount of packets

Patent Assignee: PEKKALA R (PEKK-I); PETTEY C J (PETT-I); SCHREPPPEL C L (SCHR-I)

Inventor: PEKKALA R; PETTEY C J; SCHREPPPEL C L

Number of Countries: 001 Number of Patents: 001

Patent Family:

Patent No	Kind	Date	Applicat No	Kind	Date	Week
US 20020085493	A1	20020704	US 2000740694	A	20001219	200270 B

Priority Applications (No Type Date): US 2000740694 A 20001219

Patent Details:

Patent No	Kind	Lan Pg	Main IPC	Filing Notes
US 20020085493	A1	27	H04L-001/00	

Abstract (Basic): US 20020085493 A1

NOVELTY - A memory with specific size is provided for buffering the packets, and the **flow control** credits for advertising are transmitted to a resource buffer device with size greater than the memory. The transmission of the packets is stopped by transmitting **flow control** credits to the device when it is determined that the memory is filled with specific amount of packets.

DETAILED DESCRIPTION - INDEPENDENT CLAIMS are included for the following:

- (1) Packet **flow control** method;
- (2) Packet buffering system;
- (3) **Infiniband** device; and
- (4) Buffering system.

USE - For buffering packets transmitted to **infiniband** port by **infiniband** device (claimed), to transfer data between computers and peripheral devices such as storage devices and network interface devices.

ADVANTAGE - Allows **infiniband** port to support more data virtual lanes while maintaining full **infiniband** link bandwidth through over-advertising of buffering resources. Enables supporting multiple **infiniband** ports with lesser memory requirement.

DESCRIPTION OF DRAWING(S) - The figure shows the flowchart explaining the operation of buffering system to perform over-advertising of buffering resources.

pp; 27 DwgNo 12/18

Title Terms: BUFFER; METHOD; PACKET; TRANSMIT; PORT; TRANSMIT; FLOW; CONTROL; CREDIT; RESOURCE; BUFFER; DEVICE; STOP; TRANSMISSION; PACKET; MEMORY; FILLED; SPECIFIC; AMOUNT; PACKET

Derwent Class: T01; W01

International Patent Class (Main): H04L-001/00

File Segment: EPI

13/5/13 (Item 7 from file: 350)
DIALOG(R) File 350:Derwent WPIX
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014374924 **Image available**
WPI Acc No: 2002-195627/200225
XRPX Acc No: N02-148651

Packet transmission scheduling method in communication network system,
involves grouping buffered packets into independent threads, based on
destination address, size or function of packets

Patent Assignee: SUN MICROSYSTEMS INC (SUNM)

Inventor: RYGH H

Number of Countries: 096 Number of Patents: 002

Patent Family:

Patent No	Kind	Date	Applicat No	Kind	Date	Week
WO 200197469	A2	20011220	WO 2001NO247	A	20010612	200225 B
AU 200182688	A	20011224	AU 200182688	A	20010612	200227

Priority Applications (No Type Date): US 2000593450 A 20000614

Patent Details:

Patent No Kind Lan Pg Main IPC Filing Notes
WO 200197469 A2 E 9 H04L-012/56

Designated States (National): AE AG AL AM AT AU AZ BA BB BG BR BY BZ CA
CH CN CO CR CU CZ DE DK DM DZ EC EE ES FI GB GD GE GH GM HR HU ID IL IN
IS JP KE KG KP KR KZ LC LK LR LS LT LU LV MA MD MG MK MN MW MX MZ NO NZ
PL PT RO RU SD SE SG SI SK SL TJ TM TR TT TZ UA UG US UZ VN YU ZA ZW
Designated States (Regional): AT BE CH CY DE DK EA ES FI FR GB GH GM GR
IE IT KE LS LU MC MW MZ NL OA PT SD SE SL SZ TR TZ UG ZW

AU 200182688 A H04L-012/56 Based on patent WO 200197469

Abstract (Basic): WO 200197469 A2

NOVELTY - The **buffered** packets for transmission are grouped into independent threads, based on destination address, size or function of the packets. A scheduling algorithm is applied to the grouped packets for selecting the subsequent packet for transmission.

DETAILED DESCRIPTION - An INDEPENDENT CLAIM is also included for packet transmission scheduler.

USE - For communication network systems with point-to-point links. Especially for use in network systems with NxN cross bar switch and virtual channel flow control .

ADVANTAGE - Provides multithreading and non-blocking packet transmission scheduling with fairness and optimal exploitation of connection resources based on the network traffic.

DESCRIPTION OF DRAWING(S) - The figure shows the point-to-point link with transmitter and receiver.

pp; 9 DwgNo 1/1

Title Terms: PACKET; TRANSMISSION; SCHEDULE; METHOD; COMMUNICATE; NETWORK; SYSTEM; GROUP; **BUFFER** ; PACKET; INDEPENDENT; THREAD; BASED; DESTINATION; ADDRESS; SIZE; FUNCTION; PACKET

Derwent Class: W01

International Patent Class (Main): H04L-012/56

File Segment: EPI

13/5/14 (Item 8 from file: 350)

DIALOG(R) File 350:Derwent WPIX

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012698712 **Image available**

WPI Acc No: 1999-504821/199942

XRPX Acc No: N99-377629

Virtual path capacitance control procedure for ATM communication - involves changing cell transmitting rate based on transmittable rate of cell

Patent Assignee: NIPPON TELEGRAPH & TELEPHONE CORP (NITE)

Number of Countries: 001 Number of Patents: 001

Patent Family:

Patent No	Kind	Date	Applicat No	Kind	Date	Week
JP 11220474	A	19990810	JP 9821911	A	19980203	199942 B

Priority Applications (No Type Date): JP 9821911 A 19980203

Patent Details:

Patent No	Kind	Lan Pg	Main IPC	Filing Notes
JP 11220474	A	11	H04L-012/28	

Abstract (Basic): JP 11220474 A

NOVELTY - An assignment band to each virtual path is calculated, in response to weight (W), which is assigned based on number of arrival call to the waiting buffer. A virtual path switching system assigns band to resource management cell, based on its transmission rate. Then, cell transmitting rate is changed based on transmittance rate of cell.

USE - For ATM communication.

ADVANTAGE - Warranty of minimum through-put of each user is provided.

Dwg.1/10

Title Terms: VIRTUAL; PATH; CAPACITANCE; CONTROL; PROCEDURE; ATM;

COMMUNICATE; CHANGE; CELL; TRANSMIT; RATE; BASED; RATE; CELL

Derwent Class: W01

International Patent Class (Main): H04L-012/28

International Patent Class (Additional): H04Q-003/00

File Segment: EPI

13/5/17 (Item 11 from file: 350)
DIALOG(R) File 350:Derwent WPIX
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012115479 **Image available**
WPI Acc No: 1998-532391/199845
XRPX Acc No: N98-415398

Determination of maximum allowable cell rate for output link for communications network - calculating maximum allowable flow rate as function of number of active flows through output link, and calculating alpha smoothing parameters for number of virtual channel

Patent Assignee: CABLETRON SYSTEMS INC (CABL-N); ENTERASYS NETWORKS INC (ENTE-N); CALDETRON SYSTEMS INC (CALD-N)

Inventor: CHARNEY A; CHARNY A

Number of Countries: 081 Number of Patents: 008

Patent Family:

Patent No	Kind	Date	Applicat No	Kind	Date	Week
WO 9843395	A1	19981001	WO 98US6199	A	19980327	199845 B
AU 9865910	A	19981020	AU 9865910	A	19980327	199909
US 5956322	A	19990921	US 97825201	A	19970327	199945
US 5978357	A	19991102	US 97826235	A	19970327	199953
EP 972382	A1	20000119	EP 98912116	A	19980327	200009
			WO 98US6199	A	19980327	
AU 717162	B	20000316	AU 9865910	A	19980327	200024
CA 2285086	C	20030812	CA 2285086	A	19980327	200360
			WO 98US6199	A	19980327	
EP 1381192	A1	20040114	EP 98912116	A	19980327	200410
			EP 200377139	A	19980327	

Priority Applications (No Type Date): US 97826235 A 19970327; US 97825201 A 19970327

Patent Details:

Patent No Kind Lan Pg Main IPC Filing Notes

WO 9843395 A1 E 32 H04L-012/56

Designated States (National): AL AM AT AU AZ BA BB BG BR BY CA CH CN CU CZ DE DK EE ES FI GB GE GH GM GW HU ID IL IS JP KE KG KP KR KZ LC LK LR LS LT LU LV MD MG MK MN MW MX NO NZ PL PT RO RU SD SE SG SI SK SL TJ TM TR TT UA UG US UZ VN YU ZW

Designated States (Regional): AT BE CH DE DK EA ES FI FR GB GH GM GR IE IT KE LS LU MC MW NL OA PT SD SE SZ UG ZW

AU 9865910 A H04L-012/56 Based on patent WO 9843395

US 5956322 A H04L-012/26

US 5978357 A G01R-031/08

EP 972382 A1 E H04L-012/56 Based on patent WO 9843395

Designated States (Regional): DE FR GB

AU 717162 B H04L-012/56 Previous Publ. patent AU 9865910

Based on patent WO 9843395

CA 2285086 C E H04L-012/56 Based on patent WO 9843395

EP 1381192 A1 E H04L-012/56 Div ex application EP 98912116

Div ex patent EP 972382

Designated States (Regional): DE FR GB

Abstract (Basic): WO 9843395 A

The method of determining a maximum allowed cell rate (MACR) for an output link for an ATM switch in a phantom flow control method, the output link having a threshold queue value and a previous MACR value. The method involves ascertaining a total bandwidth associated with the output link. A portion of the total bandwidth is reserved, and a reserved portion of the total bandwidth is subtracted from the total bandwidth for the output link to produce an adjusted free bandwidth.

An unused link capacity of the adjusted free bandwidth is compared to the previous MACR value. Based on the comparison, the MACR value is calculated. The calculation step involves estimating a number of virtual channels flowing into the output link over a selected time interval to produce a channel count, and ascertaining an alpha value based on the channel count.

ADVANTAGE - Improves stability of phantom flow control, and enables operation without network switches such as shared memory

switches . Prevents portion of total bandwidth of output link from being used when maximum allowable cell rate through output link is calculated.

Dwg.3/8

Title Terms: DETERMINE; MAXIMUM; ALLOW; CELL; RATE; OUTPUT; LINK; COMMUNICATE; NETWORK; CALCULATE; MAXIMUM; ALLOW; FLOW; RATE; FUNCTION; NUMBER; ACTIVE; FLOW; THROUGH; OUTPUT; LINK; CALCULATE; ALPHA; SMOOTH; PARAMETER; NUMBER; VIRTUAL; CHANNEL

Derwent Class: W01

International Patent Class (Main): G01R-031/08; H04L-012/26; H04L-012/56

International Patent Class (Additional): G06F-011/00; G08C-015/00;

H04Q-011/04

File Segment: EPI

13/5/23 (Item 17 from file: 350)
DIALOG(R) File 350:Derwent WPIX
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010718390 **Image available**
WPI Acc No: 1996-215345/199622
XRPX Acc No: N96-180770

Virtual path capacitive management device in communication networks -
has control unit which controls virtual path capacity so as to
enable tracking of judgement result, based on measured number of cells
storing bit of information

Patent Assignee: NIPPON TELEGRAPH & TELEPHONE CORP (NITE)

Number of Countries: 001 Number of Patents: 001

Patent Family:

Patent No	Kind	Date	Applicat No	Kind	Date	Week
JP 8079289	A	19960322	JP 94207668	A	19940831	199622 B

Priority Applications (No Type Date): JP 94207668 A 19940831

Patent Details:

Patent No	Kind	Lan Pg	Main IPC	Filing Notes
JP 8079289	A	7	H04L-012/42	

Abstract (Basic): JP 8079289 A

The device (1) manages the information content communicated through virtual paths between multiple nodes in communication network (100). Units of information referred as self travel between nodes through bus mutually. The communication is performed by multiplex control of information through virtual paths in each node by adapting media access control protocol.

The DQDB is used, whenever the cell reaches the terminal point of virtual path and the number of cells storing bit of information are then measured. Based on the measured result, excess and deficiency of virtual path capacity is judged. A control unit controls the virtual path capacity so as to enable tracking of judgement result.

ADVANTAGE - Reduces rate of invalid cell during data transfer.
Improves control efficiency. Lowers communication cost.

Dwg.1/6

Title Terms: VIRTUAL; PATH; CAPACITANCE; MANAGEMENT; DEVICE; COMMUNICATE;
NETWORK; CONTROL; UNIT; CONTROL; VIRTUAL; PATH; CAPACITY; SO; ENABLE;
TRACK; JUDGEMENT; RESULT; BASED; MEASURE; NUMBER; CELL; STORAGE; BIT;
INFORMATION

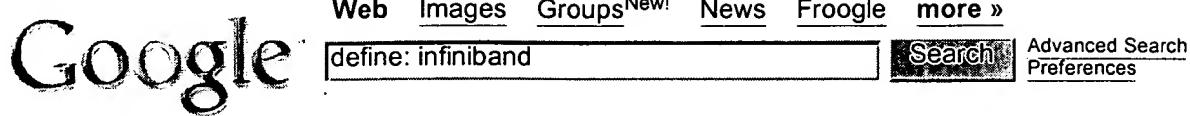
Index Terms/Additional Words: DISTRIBUTED; QUEUE ; DUAL; BUS

Derwent Class: W01

International Patent Class (Main): H04L-012/42

International Patent Class (Additional): H04L-012/28; H04L-012/46;
H04Q-003/00

File Segment: EPI



Web

Definitions of Infiniband on the Web:

InfiniBand Architecture is an industry standard, channel-based, switched fabric, interconnect architecture for servers. InfiniBand architecture changes the way servers are built, deployed, and managed.

www.mathstar.com/Technology_Glossary_IJKL.htm

A high speed short-range interconnect designed as a replacement for PCI

www.sagitta-ps.com/support/sagitta_glossary.htm

– Derived from “infinite bandwidth.” A switched-fabric I/O technology that ties together servers, storage devices and network devices, easing the bottlenecks created by data-intensive files such as streaming video, voice and audio.

www.iomega.com/europe/support/english/documents/11240e.html

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Ref #	Hits	Search Query	DBs	Default Operator	Plurals	Time Stamp
L1	2	"20020159385".did.	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2005/01/28 15:11
L2	2	"20020085493".did.	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2005/01/28 15:11
L3	56	(resource or bandwidth or data adj rate) with management same shar\$4 same (ressource or bandwidth) same (congest\$4 or load\$4 adj balanc\$4)	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2005/01/28 15:12
L4	4	(control\$4 same flow\$5) same infiniband with architecture same (bandwidth or data adj rate or resource)	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2005/01/28 15:17
L5	0	709/4\$.ccls. and (control\$4 same flow\$5) same infiniband same (bandwidth or data adj rate or resource)	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2005/01/28 15:18
L6	12	(control\$4 same flow\$5) same infiniband same (bandwidth or data adj rate or resource)	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2005/01/28 15:18

L7	106	((("5898670") or ("6108304") or ("6370475") or ("6405132") or ("4420289") or ("4984264") or ("4914959") or ("5872949") or ("6046979") or ("4251025") or ("4530805") or ("4880041") or ("5520160") or ("5839396") or ("5940372") or ("6112267") or ("6175554") or ("6192406") or ("4258424") or ("4270347") or ("4284942") or ("4313465") or ("4329120") or ("4335696") or ("4360132") or ("4440378") or ("4456223") or ("4539108") or ("4580620") or ("4598541") or ("4769810") or ("4769811") or ("4781536") or ("4809730") or ("4813858") or ("4880376") or ("4889599") or ("4944676") or ("5014265") or ("5024200") or ("5205359") or ("5315586") or ("5363830") or ("5435188") or ("5540205") or ("5544195") or ("5591923") or ("5596603") or ("5675576") or ("5748901")).PN.	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2005/01/28 15:29
L8	38	control\$4 same flow\$4 same (HCA or host adj channel adj adapter) and infiniband and virtual adj lane	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2005/01/28 15:36
L9	12	infiniband and virtual adj lane same (bandwidth or capacity or data adj rate)	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2005/01/28 15:53
L10	91	control\$4 same flow\$4 same (HCA or host adj channel adj adapter)	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2005/01/28 16:00

Ref #	Hits	Search Query	DBs	Default Operator	Plurals	Time Stamp
L1	2	"20020159385".did.	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2005/01/28 15:11
L2	2	"20020085493".did.	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2005/01/28 15:11
L3	56	(resource or bandwidth or data adj rate) with management same shar\$4 same (ressource or bandwidth) same (congest\$4 or load\$4 adj balanc\$4)	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2005/01/28 15:12
L4	4	(control\$4 same flow\$5) same infiniband with architecture same (bandwidth or data adj rate or resource)	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2005/01/28 16:35
L5	0	709/4\$.ccls. and (control\$4 same flow\$5) same infiniband same (bandwidth or data adj rate or resource)	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2005/01/28 15:18
L6	12	(control\$4 same flow\$5) same infiniband same (bandwidth or data adj rate or resource)	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2005/01/28 15:18

L7	106	((("5898670") or ("6108304") or ("6370475") or ("6405132") or ("4420289") or ("4984264") or ("4914959") or ("5872949") or ("6046979") or ("4251025") or ("4530805") or ("4880041") or ("5520160") or ("5839396") or ("5940372") or ("6112267") or ("6175554") or ("6192406") or ("4258424") or ("4270347") or ("4284942") or ("4313465") or ("4329120") or ("4335696") or ("4360132") or ("4440378") or ("4456223") or ("4539108") or ("4580620") or ("4598541") or ("4769810") or ("4769811") or ("4781536") or ("4809730") or ("4813858") or ("4880376") or ("4889599") or ("4944676") or ("5014265") or ("5024200") or ("5205359") or ("5315586") or ("5363830") or ("5435188") or ("5540205") or ("5544195") or ("5591923") or ("5596603") or ("5675576") or ("5748901")).PN.	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2005/01/28 15:29
L8	38	control\$4 same flow\$4 same (HCA or host adj channel adj adapter) and infiniband and virtual adj lane	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2005/01/28 15:36
L9	12	infiniband and virtual adj lane same (bandwidth or capacity or data adj rate)	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2005/01/28 15:53
L10	91	control\$4 same flow\$4 same (HCA or host adj channel adj adapter)	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2005/01/28 16:00
L11	0	(control\$4 same flow\$5) same infiniband with architecture same (bandwidth or data adj rate or resource) and overflow	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2005/01/28 16:37
L12	0	(control\$4 same flow\$5) same infiniband with architecture same (bandwidth or data adj rate or resource) and over near flow	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2005/01/28 17:06

L13	1	(control\$4 same flow\$5) same infiniband same (bandwidth or data adj rate or resource) and over near flow	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2005/01/28 17:07
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